

## Exercise 2 – Introduction to Cryptography 89-656

Due Date: November 11, 2018

**Exercise 1:** The best algorithm known today for finding the prime factors of an  $n$ -bit number runs in time  $2^{c \cdot n^{\frac{1}{3}} (\log n)^{\frac{2}{3}}}$ . Assuming 4Ghz computers and  $c = 1$  (and that the units of the given expression are clock cycles), estimate the size of numbers that cannot be factored for the next 100 years.

**Exercise 2:** Prove formally that the definition of indistinguishable encryptions in the presence of an eavesdropping adversary) cannot be satisfied if  $\Pi$  can encrypt arbitrary-length messages and the adversary is *not* restricted to output equal-length messages in experiment  $\text{PrivK}_{\mathcal{A}, \Pi}^{\text{eav}}$ .

**Exercise 3:** Let  $G$  be a pseudorandom generator where  $|G(s)| > 2 \cdot |s|$ .

1. Define  $G'(s) \stackrel{\text{def}}{=} G(s0^{|s|})$ . Is  $G'$  necessarily a pseudorandom generator?
2. Define  $G'(s) \stackrel{\text{def}}{=} s_1, G(s_2)$ , where  $s = s_1, s_2$  and  $s_1, s_2 \in \{0, 1\}^{|s|/2}$ . Is  $G'$  necessarily a pseudorandom generator?

**Exercise 4:** Let  $F$  be a length-preserving pseudorandom function. For the following constructions of a keyed function  $F' : \{0, 1\}^n \times \{0, 1\}^{n-1} \rightarrow \{0, 1\}^{2n}$ , state whether  $F'$  is a pseudorandom function. If yes, prove it; if not, show an attack.

1.  $F'_k(x) \stackrel{\text{def}}{=} F_k(0||x) || F_k(1||x)$ .
2.  $F'_k(x) \stackrel{\text{def}}{=} F_k(0||x) || F_k(x||1)$ .

**Exercise 5:** Consider a variant of CBC-mode encryption where the sender simply increments the  $IV$  by 1 each time a message is encrypted (rather than choosing  $IV$  at random each time); this is called *nonce-based* encryption. Show that the resulting scheme is *not* CPA-secure.

**Exercise 6:** Let  $F$  be a pseudorandom permutation. Consider the mode of operation in which a uniform value  $\text{ctr} \in \{0, 1\}^n$  is chosen, and the  $i$ th ciphertext block  $c_i$  is computed as  $c_i := F_k(\text{ctr} + i + m_i)$ . Show that this scheme does not have indistinguishable encryptions in the presence of an eavesdropper.